Combining Statistical Intervals and Market Prices: The Worst Case State Price Distribution

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Some key words and phrases: Incompleteness, Statistical uncertainty, Value at risk.

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ABSTRACT

The paper shows how to combine (historical) statistical data and (current) market prices to form regions and trading strategies. This gives rise to a “worst case” state price distribution, which provides sharp price bounds for all convex European options.

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1. Introduction. The pricing and hedging of options usually presupposes a known probability distribution $P$ for the price $S$ of the underlying security. When $P$ is not known, one approach is to find a prediction set for, say, the volatility of $S$, and then to hedge in such a way that the option liability is covered whenever the prediction set is realized (Avellaneda, Levy and Paras (1995), Lyons (1995), Mykland (2000, 2003a)). This procedure, however, fails to take account of the values of market traded options on the same security. This paper will show in the context of convex European options that such values can be incorporated in a uniform manner with the help of what we term a worst case distribution. The development is related to earlier work by Bergman (1996), Frey and Sin (1999), Frey (2000) and Mykland (2003b).

To describe our results in this paper, we begin with the cast:

(i) The securities that are traded in the market:
- $S = (S_t)_{0 \leq t \leq T}$, the price process of a stock that pays no dividend.
- $\Lambda = (\Lambda_t)_{0 \leq t \leq T}$, the price of a zero coupon bond maturing at $T$, with value one dollar (or euro, or yuen, or krone).
- European call and put options maturing at $T$ (see Section 2.1).

We can think of the value $S_t^* = S_t/\Lambda_t$ as the price of a forward contract on the stock $S$ with maturity $T$. We shall assume that $S^*$ is governed by an unknown probability $P$ which belongs to a class $Q$ of distributions. The main requirement on $P$ is that $S^*$ be an Ito process

$$dS_t^* = \mu_t S_t^* dt + \sigma_t S_t^* dW_t^*$$

(1.1)

where $(\mu_t)$ and $(\sigma_t)$ are random processes and $(W_t^*)$ is a Wiener process.

(ii) A prediction bound on the volatility $\sigma_t^2$, in the form of a prediction interval $I^{\Xi^+}$

$$I^{\Xi^+} = \{(\sigma_t) : \int_0^T \sigma_u^2 du \leq \Xi^+\}.$$  

(1.2)

(iii) A European payoff $g(S_T)$ to be made at time $T$, where $g$ will mostly be taken to be convex.

The problem we wish to solve is the following. We look for a process $(V_t)$ with two properties. $V_t$ must be the value of a self financing dynamic portfolio in the market traded securities (see Section 2.2 for the definition of this concept). Also, $V_T$ must cover the option liability if the prediction set
is realized, i.e.,

\[ V_T \geq g(S_T) \quad P - \text{a.s. on } I^\Xi^+, \text{ for all } P \in \mathcal{Q}. \]  

(1.2a)

In particular, we wish to find the amount \( V_0(g) \) which is the smallest starting value for such a self financing portfolio. A similar setup involving more general prediction sets and market traded securities is given in Mykland (2003a), which discusses the relevant concepts in some detail.

What is special about the development in the current paper is that we show the existence of a mapping from the prediction bound \( \Xi^+ \) to a cumulative distribution \( F^{\Xi^+} \) on \( S_T \):

\[ \Xi^+ \rightarrow F^{\Xi^+} \]  

(1.3)

so that for all (non strictly) convex \( g \) which do not grow too fast, the conservative starting value for prediction set \( I^{\Xi^+} \) and payoff \( g(S_T) \) at \( T \) is

\[ V_0(g) = \Lambda_0 \int g(s) dF^{\Xi^+}(s). \]  

(1.4)

In other words: there is one \( F \) which can be taken as the worst case state price distribution for all convex payoffs \( g(S_T) \). Convex includes calls, puts.

\( F^{\Xi^+} \) is what we call the worst case distribution for the market structure and prediction set described above. In consequence, since \( F^{\Xi^+} \) is independent of the payoff function \( g \), one does not need to compute the value \( V_0(g) \) for each \( g \), but instead can find the distribution \( F^{\Xi^+} \). Also, a distribution function is a more conceptual object.

2. Description and main theorem.

2.1. The worst case distribution. For reasons of mathematical convenience, assume that all market traded options are European puts. This is no restriction on results, as puts and calls can be converted to each other via put-call parity, see Chapter 7.4 (pp. 174-175) of Hull (1999). In practice, one would want to take the most liquid of the put and the call to minimize the transaction cost. A put option with strike price \( K \) has value \( (K - S_T)^+ \) at maturity \( T \). Its market price at time \( t \) will be denoted by \( P^K_t \), and for the discounted quantity, we use \( P^{K^*}_t = P^K_t / \Lambda_t \). A discussion of a formulation in terms of calls is given in Section 7.
We suppose that the market traded puts have strike prices \( K_1, ..., K_q \).

To describe \( F^{\Xi^+} \), introduce the standardized process
\[
d\tilde{S}_t = \tilde{S}_t d\tilde{W}_t, \quad \tilde{S}_0 = S_0^*(= S_0/\Lambda_0)
\]
where \( \tilde{W} \) is a standard Brownian motion. The corresponding probability distribution will be called \( \tilde{P} \). Then
\[
F^{\Xi^+}(s) = \tilde{P}(\tilde{S}_{\tau \wedge \Xi^+} \leq s)
\]
where:
\[
\tau = \inf\{t : t \geq t_i \text{ and } \tilde{S}_r = K_i \text{ for some } i, \ 1 \leq i \leq q\}
\]
and where the \( t_1, ..., t_q \) are nonrandom, independent of \( \Xi^+ \), and the solution of
\[
\tilde{E}(K_i - \tilde{S}_\tau)^+ = P_0^{K_i}/\Lambda_0 \quad 1 \leq i \leq q
\]
where the left hand side of (2.4) is taken as a function of the \( t_i \)s through (2.3).

**Proposition 1.** If (2.3)-(2.4) has a solution, it is unique.

This follows directly from Theorem 2 in Section 5 below. The case of no solution will be considered in Section 3.2. We shall suppose that the solution of (2.4) satisfies
\[
\max_{i=1, ..., q} t_i \leq \Xi^+, \quad (2.5)
\]
without which there is “arbitrage” in a sense to be discussed also in Section 3.2.

**2.2. Trading strategies: Theoretical considerations.** We define the class \( Q \) given the initial discounted values \( S_0^* \) and \( P_0^{K_i^*} \), \( 1 \leq i \leq q \), as follows:

**Definition.** \( Q \) is a collection of distributions on the set of functions \( \Omega = \mathcal{C}[0, T] \times \mathcal{D}[0, T]^q \), and \( (S^*_t, P_t^{K_1^*}, ..., P_t^{K_q^*}) \) is the coordinate process. Every \( P \in Q \) must satisfy (1.1), and the coordinate process must have the correct initial value \( (S_0^*, P_0^{K_1^*}, ..., P_0^{K_q^*}) \) with probability one. Also, for given \( P \), the process \( (\sigma_t) \) must be bounded \( P \)-a.s. (but one does not need to know the bound). Finally, each \( P \) must be mutually absolutely continuous with a \( P^* \) under which the coordinate process is a martingale. The collection of such \( P^* \)s will be called \( Q^* \).
The requirement of equivalence to a “risk neutral measure” $P^*$ is the most convenient way to avoid arbitrage opportunities in the market. Not only is $S^*$ a martingale under $P^*$, but processes $P^{K_i}$ can be found, with correct initial value $P^*_0$, such that $P^{K_i*}$ is a martingale.

The filtration describing the market will be called $(\mathcal{F}_t)$ and can be any that is right continuous, and which makes the coordinate process adapted and martingales under all $P^* \in \mathcal{Q}^*$. This $(\mathcal{F}_t)$ can be either the smallest filtration with these properties, or anything bigger satisfying the same criteria, such as the “analytic completion” discussed in Mykland (2003a). The latter is particularly useful from a statistical perspective. All processes are taken to be càdlàg and adapted to $(\mathcal{F}_t)$.

A self financing dynamic portfolio $(V_t)$, with discounted value $V^*_t = V_t/\Lambda_t$, is defined as a process which, for any $P \in \mathcal{Q}$, can be represented by $V^* = H^* - D^*$, where $D^*$ is non decreasing (“dividend”), and $H^*$ is a stochastic integral with respect to $S^*$ and the $P^{K_i*}$. Stochastic integrals are as defined in Chapter I.4d (pp. 46-51) of Jacod and Shiryaev (2003). The random variables $\{H^*_{\lambda}\}$ must be uniformly integrable for all $P^* \in \mathcal{Q}^*$, where the $\lambda$ describe the set of stopping times taking values in $[0,T]$.

We can confine ourselves to considering discounted quantities by virtue of numeraire invariance, see Chapter 6 of Duffie (1996). In the current situation where we discount by the zero coupon bond $\Lambda$, $\mathcal{F}^\Xi$ and the option liabilities $g(S_T)$ and $(K - S_T)^+$ also only depend on discounted quantities, since $S_T = S^*_T$. All conditions, therefore, can and are directly imposed on the discounted processes. The author learned this device from the paper of El Karoui, Jeanblanc-Piqué, and Shreve (1998).

2.3. **Form of the self financing portfolio.** The trading strategy that starts with $V_0(g)$ only requires a static position in the market traded options. We explain this in the following.

At the outset of trading, suppose one takes a position of $\lambda_i$ units in the put option with strike price $K_i$. We let this position be static, in the sense that we hold it without change until expiry at time $T$. The problem then changes to that of covering a liability of the form $h_\lambda(S_T)$, where

$$ h_\lambda(s) = g(s) - \sum_{i=1}^{q} \lambda_i[(K_i - s)^+ - P_0^{K_i}] $$

(2.6)

This is since a loan of $P_0^K$ dollars at time 0 requires a repayment of $P_0^{K^*}$ dollars at maturity.

The two problems are equivalent, and $V_0(h_\lambda) = V_0(g)$. It turns out, however, that there is one value of $\lambda = (\lambda_1, ..., \lambda_q)$ so that the dynamic hedge for liability $h_\lambda(S_T)$ only involves trading in the forward contract $S^*$ (in other words, $H^*$ is a stochastic integral over $S^*$ only).

Practically, this is important because transaction costs are normally higher in the market traded options than in the forward contract $S^*$ (which, if need be, can be created by securities $S$ and $\Lambda$). We otherwise ignore the issue of trading cost in this paper.

2.4. **The main result.**

**Theorem 1.** Assume that (2.3)-(2.4) has a solution. Also suppose (2.5). Then there exists a mapping (1.3) so that for all (non strictly) convex $g$ satisfying $|g(s)| \leq a + bs^{2-\epsilon}$ for some $\epsilon > 0$, the conservative starting value $V_0(g)$ for prediction set $F^{\Xi^+}$ and payoff $g(S_T)$ at $T$ is given by (1.4). $F^{\Xi^+}$ is given by (2.2) above. There is a trading strategy that starts with value $V_0(g)$, and only requires a static position in the market traded options, as described in Section 2.3 above.

The theorem is proved in the Appendix. – The general issue of how to compute the worst case distribution and its corresponding hedging strategy is discussed below in Section 6. First, however, an interpretation of the $t_i$s.

3. **Implied volatilities, and arbitrage.** Note first that the construction in Section 2.1 is a conversion of time to volatility scale. The $t_i$’s can be seen as a form of implied volatilities. First, consider the case where this is exactly true.
3.1. Connection to implied volatility. The Black-Scholes (1973)-Merton (1973) form of the price at time 0 of a European put option with strike $K$ is $BSMP(S_0, -\log \Lambda_0, \sigma^2 T)$, where

$$BSMP(S, R, \Xi) = K \exp(-R) \Phi(-d_2) - S \Phi(-d_1)$$

with

$$d_{1,2} = (\log(S/K) + R \pm \Xi/2) / \sqrt{\Xi},$$

and where the instantaneous volatility $\sigma_t^2$ is taken to be constant and equal to $\sigma^2$. Also, $\Phi$ is the standard normal cumulative distribution function.

Of course, the model underlying this formula may not be valid, and prices do not generally behave as if it were, but it is customary to invert it to find so-called implied volatilities. We shall here do this on the cumulative scale.

**Definition.** The (cumulative) implied volatility at time zero for strike price $K$, $\Xi_K$, is defined by

$$BSMP(S_0, -\log \Lambda_0, \Xi_K) = P_K^0$$

It is also natural to call the $t_i$ from Theorem 1 the conditional (cumulative) implied volatility at time zero for strike price $K_i$. For both objects, we omit “cumulative” unless this is not clear from the context.

The first connection to $F^{\Xi^+}$ is now as follows.

**Example 1.** If all traded options have the same implied volatility:

$$\Xi_{K_1} = \ldots = \Xi_{K_q},$$

then

$$t_1 = \ldots = t_q = \Xi_{K_1}$$

This is easily seen from Theorem 1. In the more general case of unequal implied volatilities, $\argmin_i(t_i) = \argmin_i(\Xi_{K_i})$ (there can be several such $i$’s, or course) and $t_i$ and $\Xi_{K_i}$ must coincide for these indices $i$. Also, one can see more generally from the convexity of the put payoff that $t_i \geq \Xi_{K_i}$. ■
3.2. The $t_i$'s and arbitrage. Arbitrage is the construction of a self financing strategy which makes a profit for some $P \in \mathcal{Q}$, and which does not lose money almost surely, for all $P \in \mathcal{Q}$. For issues related to the avoidance of doubling strategies, see Chapter 6 of Duffie (1996).

There are two ways that arbitrage can occur in our setting. One is if (2.5) is violated. The other is if the system (2.3)-(2.4) has no solution. The latter case is the most clear cut (with proof in the Appendix):

**Proposition 2.** If (2.3)-(2.4) has no solution, then there is arbitrage.

The former case is one of “statistical arbitrage”, in the sense that the prediction interval $I^{\Xi^+}$ has to be realized, otherwise a loss can occur.

**Proposition 3.** Assume that (2.3)-(2.4) has a solution, but that $\max_{i=1, \ldots, q} t_i > \Xi^+$. Then there is a trading strategy which, provided $I^{\Xi^+}$ is realized, yields a positive return of at least $c^*$ at time $T$, almost surely for any $P \in \mathcal{Q}$. Here $c^*$ is a positive constant independent of $P$.

**Proof of Proposition 3.** Let $S$ be the indices for which the $t_i$s are smaller than $\Xi^+$, and let $I$ be the index $i$ corresponding to the smallest $t_i$ strictly greater than $\Xi^+$. We shall be interested in earning money on the put payoff $g(s) = (K_I - s)^+$. The following argument remains valid if $S$ is empty.

Note that the $t_i$, $i \in S$, solve (2.3)-(2.4) for this index set. Let $V_0(g)$ be the price given by Theorem 1 based on hedging in $S$, $\Lambda$ and the puts with strike $K_i$, $i \in S$. Our claim is now that

$$V_0(g) < P_0^{K_I}.$$  \hfill (3.1)

Thus, one can sell the option with payoff $g(S_T)$, start a self financing trading strategy with initial value $V_0(g)$, and be sure that the liability is covered so long as $I^{\Xi^+}$ is realized. The profit is at least $c = P_0^{K_I} - V_0(g)$, if taken at time 0, or $c^* = c/\Lambda_0$ if taken at time $T$ (any random additional profit has, of course, to be taken at time $T$).

To see (3.1), let $\tau$ be given on the form (2.3) based on $t_i$, $i \in S \cup \{I\}$. Also let $K_- = \sup\{K_i < K_I, \ i \in S\}$ and $K_+ = \inf\{K_i > K_I, \ i \in S\}$. Let $C = \{\tau \geq \Xi^+ \text{ and } S_{\Xi^+} \in (K_-, K_+)\}$. It then follows that

$$P_0^{K_I} - V_0(g) = \mathbb{E}[g(S_{\tau}) - g(S_{\tau \wedge \Xi^+})] = \mathbb{E}[X],$$
where $X$ is zero outside $C$, and otherwise positive.

4. **An implementation with data.** We give an example of how the worst case distribution can look. Table 1 provides (Bayesian) posterior quantiles for the square root of $\Xi = \int_0^T \sigma^2_t dt$ for the S&P 500. The posterior distribution is based on work by Jacquier, Polson and Rossi (1994), which analyses (among other series) the S&P 500 data recorded daily, and has also been used by Mykland (2003a), which provides further discussion.

<table>
<thead>
<tr>
<th>Posterior Coverage</th>
<th>50%</th>
<th>80%</th>
<th>90%</th>
<th>95%</th>
<th>99%</th>
</tr>
</thead>
<tbody>
<tr>
<td>Upper end $\sqrt{\Xi T}$</td>
<td>.168</td>
<td>.187</td>
<td>.202</td>
<td>.217</td>
<td>.257</td>
</tr>
</tbody>
</table>

*The posterior is conditional on $\log(\sigma_0^2)$ taking the value of the long run mean of $\log(\sigma^2)$.*

We now assume that market traded options are available for a certain number of strike prices. For a given set of securities prices, a worst case distribution given in Fig. 1, corresponding to the 90% posterior quantile. A comparison of the worst case distribution for different quantiles is given in Fig. 2.
Fig 1. Worst case distribution based on the 90% posterior quantile, and incorporating traded options with strike prices 60, 80, 100, 120, and 140. The traded options were taken to have Black-Scholes implied volatility (see Section 3) equal to the 80% posterior quantile.
Fig 2. Worst case distributions for the 80% (yellow curve), 90% (red), 95% (blue), 99% (green) and 100% (black) posterior quantiles. Traded options and their values are as in Fig. 1. Note that at each of the strike prices \( K \), more point mass is placed on \( K \) the higher the posterior coverage of the statistical interval.
5. Finding \( t_1, ..., t_q \): the general case. We here present algorithms for finding the conditional implied volatilities. Set \( P^K_t = P^K_t / \Lambda_t \). In other words, this is the discounted put price at \( t \). We start by characterizing the output, and the algorithms are stated just after the Theorem.

**Theorem 2.**

(a) Algorithms 1 and 2 yield the same result.

(b) \((2.3)-(2.4)\) has a solution if and only if either algorithm does not return a “no solution” message.

(c) In the absence of a “no solution” message, the output of either algorithm is unique, and satisfies \((2.3)-(2.4)\).

(d) If \((2.3)-(2.4)\) has a solution, then it is unique, and is provided by either algorithm.

Obviously, one of statements (b) and (d) is redundant given (c), but it seemed to improve readability to have them both there. The result is proved in the Appendix. Recall that the case of no solution to \((2.3)-(2.4)\) has been discussed in Proposition 2 above.

Before going to the two main algorithms, we give the core component of both as Algorithm 0. At the end of the section, we show some more detail on how to compute Step 1 in the following.

**Algorithm 0.** We use an index set \( \mathcal{S} \subseteq \{1, ..., q\} \), where index \( i \) corresponds to \( K_i, P^0_{K_i} \), and \( t_i \). \( t_j \)'s where \( j \) is not in \( \mathcal{S} \) have either already been found, or are irrelevant. Also, a provisional version of \( \tau \) is given.

0. If, for any \( i \in \mathcal{S} \), \( \bar{E}(K_i - \bar{S}_\tau) + \leq P^0_{K_i} \), then the algorithm terminates with a “no solution” message. Otherwise:

1. For \( i \in \mathcal{S} \), find \( t_i \) by:

\[
\bar{E}(K_i - \bar{S}_{\tau \wedge t_i})^+ = P^0_{K_i}.
\]

2. Remove all \( i \) corresponding to the smallest \( t_i \) from \( \mathcal{S} \) (there can be ties between the \( i \)s).

3. Set

\[
\tau = \inf\{t : t \geq t_i \text{ and } \bar{S}_\tau = K_i \text{ for } i \text{ not in } \mathcal{S}\}.
\]
The algorithms described in Theorem 2 are then given by

**ALGORITHM 1.** Finds $t_1, ..., t_q$ in accordance with Theorem 2. This is the loop version. If any option value is non-positive, the algorithm terminates with a “no solution” message. Otherwise, set initial values: $S = \{1, ..., q\}$ and $\tau = +\infty$. Then

Loop:

1. Go through Steps 0-3 in Algorithm 0
2. Repeat loop until $S$ is empty (unless Step 0 has triggered early termination)

For aficionados of recursion, an alternative scheme would be the following.

**ALGORITHM 2.** Finds $t_1, ..., t_q$ in accordance with Theorem 2. This is the recursive version.

For index sets $S \subseteq \{1, ..., q\}$, and for stopping times $\tau$, a functional $F = F(S, \tau)$ is defined below. It returns either a solution, or detects is absence.

If any option value is non-positive, the algorithm terminates with a “no solution” message. Otherwise, the overall solution to the algorithm is $F(\{1, ..., q\}, \tau = +\infty)$.

**Definition of $F$:**

1. If $S$ is the empty set, then $F(S, \tau)$ returns no information. Otherwise
2. Carry out item 0-3 from Algorithm 0. Let $i_1 < ... < i_r$ be the indices picked out by Step 2. Then
3. Define $S_1 = S \cap \{1, ..., i_1 - 1\}$, $S_{r+1} = S \cap \{i_r + 1, ..., q\}$, and, for $v = 2, ..., r$, $S_v = S \cap \{i_{v-1} + 1, ..., i_v - 1\}$
4. $F(S_v, \tau)$ returns the $\{t_j, \ j \in S_v\}$ for all $v$. Hence $F(S, \tau)$ returns $\{t_j, \ j \in S\}$. If either of the $r + 1$ recursions returns a “no solution” message, then $F(S, \tau)$ returns a “no solution” message.

A useful fact is the following.

**Proposition 4.** In Algorithm 1, if index $i$ is picked out from $S$ in an earlier pass through the loop than $j$, then $t_i < t_j$. Similarly, in Algorithm 2, if index $i$ is picked out from $S$ at an earlier point in the recursion than $j$, then $t_i < t_j$. 
Remark 1. The calculation of $t_i$ in Step 1 in Algorithm 0 can be implemented as follows. Let $t_-$ be the largest previous value of $t_j$'s selected by Step 3 earlier in the loop or the recursion, or set $t_- = 0$ if none has been selected. In view of Step 3 and of Proposition 4, and also of the requirement that traded options values be positive (so $\Xi_{k_i} > 0$), $t_i > t_-$. Let $a = \max\{j \not\in S, j < i\}$ or $a = 0$ if this set is empty. Similarly, let $b = \min\{j \not\in S, j > i\}$ or $b = q + 1$ if this set is empty. Take $K_0 = 0$ and $K_{q+1} = +\infty$. Note that, on the set $A = \{K_a \leq \bar{S}_t \leq K_b\} \cap \{\tau \geq t_-\}$, the preexisting $\tau$ is given by

$$\tau = \inf\{t \geq t_- : \bar{S}_t = K_a \text{ or } K_b\}.$$ 

It follows that $t_i$ is given by by Step 1 via

$$\bar{E}(DO(S_{t_-}, t_-, t_i)I_A) = P_0^{K_i} - \bar{E}(K_i - \bar{S}_\tau)^+I_A$$

where, for $K_a \leq s \leq K_b$,

$$DO(s, t_-, t) = \bar{E}[(K_i - \bar{S}_\tau)^+|S_{t_-} = s]$$

is the value at $t_-$ of the double barrier put option with starting value $s$ and cumulative volatility $t - t_-$ (from $t_-$ onward). Since this function is known to be continuous and strictly increasing in the volatility, it follows that (5.1) has a unique solution on $(t_-, +\infty)$ unless the algorithm has been terminated in Step 0.

6. Finding the static hedge coefficients $\lambda_i$. A remaining issue is to define the $\lambda_i$'s from (2.6). For given payoff function $g$, let $\mu$ be the second derivative measure associated with $g$. In other words, $\mu$ is a measure which satisfies $\mu([x, y]) = g'(y) - g'(x)$ for all $x, y$ outside a countable set on the real line. $\mu$ exists since $g$ is convex, see Karatzas and Shreve (1991), pp. 212-214, for details.

For any $\tau$ on the form (2.3),

$$\bar{E}g(\bar{S}_{\tau\wedge\Xi^+}) = \int \bar{E}(K - \bar{S}_{\tau\wedge\Xi^+})^+\mu(dK).$$

By the same argument as in Remark 1, $\bar{E}(K - \bar{S}_{\tau\wedge\Xi^+})^+$ is continuously differentiable in each $t_i$, and therefore so is $\bar{E}g(\bar{S}_{\tau\wedge\Xi^+})$. As a consequence of the proof of Lemma 2 in the Appendix, one
thus obtains the $\lambda_i$s through

$$\frac{\partial E g(\tilde{S}_{\tau\wedge \Xi}^+)}{\partial t_i} = \sum_{j=1}^q \lambda_j \frac{\partial E (K_j - \tilde{S}_{\tau})^+}{\partial t_i}, \quad i = 1, ..., q, \tag{6.2}$$

when this expression is evaluated at the values $t_i$ which solves equations (2.3)-(2.4). Note, from Remark 1, that $\partial E (K_j - \tilde{S}_{\tau})^+/\partial t_i$ is nonzero only when $t_j > t_i$. Hence, if one orders the indices so that $t_1 \leq ... \leq t_q$, the system of equations (6.2) involves a triangular matrix, and hence has a unique solution.

7. Other issues.

Formulation in terms of call options. The results can equally well be put in terms of call options, but one then needs an additional caveat. We use in the proof of Theorem 1 that $(K_i - \tilde{S}_{\tau\wedge t})^+$ is a uniformly integrable martingale from $\max t_i$ onward. For $(\tilde{S}_{\tau\wedge t} - K_i)^+$ the same statement is true for $i < q$, but for $i = q$ it is not: $(\tilde{S}_{\tau\wedge t} - K_q)^+$ is a martingale, but its limit at $t \to +\infty$ is zero. This can be remedied by requiring $\tau$ in (2.3) to be bounded by a constant $c$, which one can take to be $\Xi^+$, or max $t_i$, or any other number greater than either. Algorithm 1 then uses the starting value $\tau = c$ rather than $\tau = +\infty$. The reason we preferred to avoid this formulation is that it does not make it clear that $\tau$ is, in fact, independent of $\Xi^+$. – This problem does not arise when evaluating the expectation of $g(\tilde{S}_{\tau\wedge \Xi}^+)$, since $\Xi^+$ is an upper bound on the stopping time.

An alternative formulation in terms of calls would be to replace (2.4) by

$$\tilde{E}(\tilde{S}_{\tau} - K_i)^+ = C_0^{K_i}/\Lambda_0,$$

with the side condition that $\tau \leq t_q$ on the set \{\(S_{t_q} \geq K_q\}\). This is again somewhat more inelegant than the formulation with puts, which is why we have stuck with Theorem 1 as it is.
APPENDIX: PROOFS OF THEOREMS

Logical sequence of proofs. For ease of reference, the proofs are given in the order of appearance of the results in the main text. The results, however, depend on each other in a different logical sequence, and should be taken to be proved in the following order. Proposition 4 is proved from scratch. Then Theorem 2 is proved using Proposition 4, Proposition 1 is a direct corollary to Theorem 2, and Theorem 1 uses Proposition 1. Proposition 3 (proved in the main text) uses Theorem 1. Proposition 2 uses both Theorem 1 and the development in Theorem 2. Finally, Lemmas 1-2 are embedded in the proof of Theorem 1.

Proof of Theorem 1.

Exit \( \Lambda \), followed by bear. As discussed in Section 2.2, we make use of the discounting by \( \Lambda_t \) to restate the problem in terms of finding a self financing strategy in \( S^* \) and the \( PK_i^* \). First of all, the option liability \( \eta = g(S_T) \) can be reexpressed as \( \eta = g(S_T^*) \), since \( \Lambda_T = 1 \). Second, by numeraire invariance (see Duffie (1996), Chapter 6), \( V_t \) is a self financing portfolio in the securities \( S, \Lambda \) and the \( PK_i \) if and only if \( V_t^* = V_t/\Lambda_t \) is a self financing strategy in the forward contracts given by \( S^* \) and the \( PK_i^* \). Since \( V_T = V_T^* \), the liability \( \eta \) will be covered by \( V^* \) if and only if it is covered by \( V^* \). – It is enough, therefore, to prove Theorem 1 as if the \( \Lambda \) process were identically equal to 1. In other words, as if uninvested cash were stored in the mattress.

The function \( g \) can be taken to be bounded below. Without loss of generality, we can assume that the \( g \) function is bounded below. This is because, by convexity, there is a constant \( c_1 \) so that \( g_1(s) = g(s) + c_1 s \) is bounded below. One can then hedge the liability \( g(S_T) \) by instead hedging \( g_1(S_T) \), and in addition take a static position of \( -c_1 \) units of security \( S \).

Reformulation of the problem. We shall consider related problems on the set \( \Omega' = \mathbb{C}[0,T] \), with coordinate process \( (S^*_t) \), and prespecified initial value \( S^*_0 \). We work with various collections of probabilities.

\( \mathcal{R}^*_\Xi \) is the set of probabilities so that \( S^* \) is a martingale satisfying (1.1), in particular,

\[
    dS^*_t = \sigma_t S^*_t dW^*_t, \tag{A.1}
\]
and so that for given \( P^* \in \mathcal{R}^*_{\Xi^+} \) the process \((\sigma_t)\) must be bounded with probability one. We also require \( P^*(I_{\Xi^+}) = 1 \) for all \( P^* \in \mathcal{R}^*_{\Xi^+} \). Also, 

\[
\mathcal{P}^*_{\Xi^+} = \{ P^* \in \mathcal{R}^*_{\Xi^+} : E^*(K_i - S^*_T)^+ = P^*_0 K_i^* \}.
\]

Note that \( \mathcal{P}^*_{\Xi^+} \) extends to \( Q^* \) on \( \Omega \). Define 

\[
\bar{V}_0(g) = \sup_{P^* \in P^*_{\Xi^+}} E^*g(S^*_T). \tag{A.2}
\]

and 

\[
\bar{V}_0(g; \lambda) = \sup_{P^* \in R^*_{\Xi^+}} E^*h_{\lambda}(S^*_T). \tag{A.3}
\]

By the Dambis (1965)/Dubins-Schwartz (1965) time change, 

\[
\bar{V}_0(g; \lambda) = \sup_{0 \leq \tau \leq \Xi^+} \bar{E}h_{\lambda}(\bar{S}_\tau), \tag{A.4}
\]

where \( \tau \) describe all stopping times in the interval \([0, \Xi^+]\). Finally set 

\[
\bar{V}_0(g) = \inf_{\lambda} V_0(g; \lambda) \tag{A.5}
\]

First, by a Lagrange argument, we get

**Solution of (A.4)-(A.5), and equality to (A.2).** Problem (A.4) can be solved using standard procedure for American options (see Karatzas (1988), Myneni (1992), and the references therein), which yield that the supremum is attained at a stopping time \( \tau_{\lambda}^* \). The American option argument makes use of the *Snell envelope* for \( h_{\lambda} \), which reenters the discussion below:

\[
\text{SE}(s, \Xi; \lambda) = \sup_{\Xi \leq \tau \leq \Xi^+} \bar{E}(h_{\lambda}(\bar{S}_\tau) \mid S_\Xi = s).
\]

By an argument similar to that of Theorem 3 of Mykland (2003b), \( \tau_{\lambda}^* = \tau_{\lambda} \wedge \Xi^+ \), where 

\[
\tau_{\lambda} = \inf \{ t : t \geq t_i^\lambda \text{ and } \bar{S}_t = K_i \text{ for some } i, 1 \leq i \leq q \}.
\]

**Lemma 1.** Suppose that \( g \) is bounded below with \( |g(s)| \leq a + bs^{2-\epsilon} \). Also assume that \((2.3)-(2.4)\) has a solution \( \tau \) satisfying \((2.5)\), Then there is such a \( \tau \) so that \( \bar{V}_0 = \bar{E}g(\bar{S}_{\tau \wedge \Xi^+}) \).
Proof of Lemma 1. Consider a sequence of \( \lambda \) so that \( \bar{V}_0(g; \lambda) \) converges to \( \bar{V}_0(g) \). Since the \( t^\lambda_i \)'s live in the compact set \([0, \Xi^+]^q\), there is a subsequence which is convergent in the \( t^\lambda_i \)'s. Call the relevant limit \( t_i \), and define \( \tau^* \) as the limit of the \( \tau^*_\lambda \) as one passes through the subsequence. By uniform integrability, and since \( \tilde{S} \) is a continuous process, \( \tilde{E}g(S^\tau_\lambda) \to \tilde{E}g(S^\tau) \) and, for \( i = 1, \ldots, q \), \( \tilde{E}(K_i - \tilde{S}^{\tau^*_\lambda})^+ \to \tilde{E}(K_i - \tilde{S}^{\tau^*})^+ \). Also, \( \tau^* \) can be taken to be on the form \( \tau \land \Xi^+ \), where \( \tau \) is on the form (2.3), again since \( \tilde{S} \) is a continuous process.

\( \tau^* \) must satisfy (2.4), otherwise the infimum in (A.5) would not be finite. By the argument just underneath the statement of Lemma 1, \( \tau^* \) can be replaced by \( \tau \) for the purpose of satisfying (2.4).

If (2.3)-(2.5) has a solution, \( V_0(g) > -\infty \). This is because, since \((K_i - \tilde{S}^{\tau \land \lambda})^+ \) is a uniformly integrable martingale from max \( t_i \) onward, the constraint (2.4) will also be satisfied if \( \tau \) is replaced by \( \tau \land \Xi^+ \), provided max \( t_i \leq \Xi^+ \). Hence also \( V_0(g) \leq \bar{V}_0(g) \). Lemma 1 then shows that \( V_0(g) = \bar{V}_0(g) \).

Connection to the worst case distribution, and the trading strategy. Now combine Lemma 1 with Proposition 1 to see that

\[
V_0(g) = \bar{V}_0(g) = \int g(s) dF^{\Xi^+}(s).
\]

Also, observe that \( V_0(g) \) must be a lower bound for the starting value \( V^*_0 \) of any self financing strategy \( V^*_t \) for which \( V^*_T \geq g(S^*_T) \) on \( I^{\Xi^+} \). This is because any strategy would have to be self financing and solvent under each \( P^* \in P^*_\Xi \), cf. the development in Mykland (2000, 2003a), and the literature on superhedging cited in the introduction to the former of these two papers, cf. also Section 2.2 in this paper.

Existence of a self financing strategy with initial value \( \bar{V}_0(g) = \bar{V}_0(g) \). Theorem 1 will have been shown if \( V^*_0 \) can be taken to be \( \bar{V}_0(g) \). To do this, observe that

Lemma 2. Under the assumptions of Lemma 1, there is a (finite) value of \( \lambda \) to that \( V_0(g) = \bar{V}_0(g; \lambda) \). This value is the unique solution of the system of equations (6.2).

For this \( \lambda \), the process \( V^*_t = SE(S^*_t, \Xi^+ - \int_0^t \sigma^2 du, \lambda) \) satisfies our requirements for a self financing strategy in \( S^* \) that is solvent on \( I^{\Xi^+} \), for all \( P^* \in P^*_\Xi \), and hence for all \( P \in Q \). This is,
again, by the Dambis/Dubins-Schwartz time change, and by the solution for the American payoff $h_{\lambda}(S_{\tau})$ under the model (2.1).

**Proof of Lemma 2.** For a given values of $\lambda_i$, the values of $t_i$ that minimize $\bar{V}_0(g;\lambda)$ must satisfy equations (6.2). Also, by the triangular matrix argument mentioned in Section 6, this system of equations defines the $\lambda_i$’s from the minimizing $t_i$’s. When one takes the limit in Lemma 1, equation (6.2) therefore remains valid, by continuity of $\partial \bar{E}g(\bar{S}_{\tau \wedge \Xi^+})/\partial t_i$ and the $\partial \bar{E}(K_j - \bar{S}_{\tau})^+/\partial t_i$.

This ends the proof of Theorem 1.

**Proof of Proposition 2.** If (2.3)-(2.4) has no solution, this means that there will be a pass through the loop in Algorithm 1 where step 0 returns a no solution message. Let $S$ be the index set at this stage, and let $I \in S$ be an index that causes the termination condition to be triggered. The arbitrage strategy is then constructed as in the proof of Proposition 3, in view of (A.10) below.

**Proof of Proposition 4.** We only show the loop case. The recursion case is similar.

Let $t_-$ be as in Remark 1. In the first pass through the loop, Step 1 sets $t_i = \Xi_{K_i}$. This value will exceed $t_- = 0$.

We now proceed by induction, assuming that we have gone through $n$ passes of the loop, $n \geq 1$. We need to show that all $t_i$ found in Step 1 strictly exceed $t_-$. If $t_i = t_-$, then, the index $i$ would have been selected in the previous pass. If $t_i < t_-$, this means that there is a set $\{j_1, ..., j_r\}$ so that the $t_{j_k}$ have already been picked out in earlier passes through the loop, and so that

$$t_i \leq t_{j_1} \leq ... \leq t_{j_r} = t_-,$$

where at least one of the inequalities is strict. Let $l$ be the pass of the loop where $t_{j_1}$ is picked out, and let $\tau'$ be the $\tau$ given by Step (3) in pass number $l - 1$. (If $l = 1$, then $\tau' = +\infty$).

Since $(K_i - \bar{S}_{\tau})^+$ is a martingale for $t \geq t_i$, it follows that

$$\bar{E}(K_i - \bar{S}_{\tau \wedge \Xi^+})^+|\mathcal{F}_{t_i} = (K_i - \bar{S}_{\tau \wedge \Xi^+})^+ = (K_i - \bar{S}_{\tau' \wedge \Xi^+})^+. $$

(A.7)
Hence, by taking unconditional expectations, and using (A.6),

$$\tilde{E}(K_i - \tilde{S}_{\tau' \wedge t_i})^+ = P_{i}^{K_i,*}. $$

It follows that index $i$ would give rise to value $t_i$ in Step 1 of the $l$'s iteration of the loop.

Let $S_n$ be the set of unselected indices at the start of iteration $n+1$, and set $S_n^c = \{1, ..., q\} - S$. There are two possibilities in (A.6): either, for some $k < r$,

$$t_i \leq t_{j_1} = ... = t_{j_k} < t_{j_{k+1}} \leq t_-,$$

(A.8)
or there is no such $k$, in which case

$$t_i < t_{j_1} = ... = t_{j_r} = t_-.$$ (A.9)

In the event of (A.8), Step 1 of pass $l$ of the loop gives the values $t_{j_1} = ... = t_{j_k}$ for indices $\{j_1, ..., j_k\}$. Also, the indices in $\{j_{k+1}, ..., j_r\} \cup (S_n^c - \{i\})$ are rejected in Step 2. Hence, at most, $\{i, j_1, ..., j_k\}$ are selected in Step 2, and possibly only $\{i\}$. In the event of (A.9), the same reasoning applies. In any case, index $i$ is picked out in iteration $l < n + 1$, the current iteration number of the loop. Hence, again, if $t_i$ were strictly smaller than $t_-$, it would already have been picked out by the loop in a previous step.

**Proof of Theorem 2.** (a) is trivial. We do the proof of the rest in steps (i)-(iii) below. (c) follows from (i) and (ii), (d) follows from (i) and (iii), (b) follows from (c) and (d) in the statement of the theorem.

(i) Algorithm 1, provided there is no termination in Step 0 at any point in the loop, provides a unique result $t_1, ..., t_q$.

This follows from Proposition 4 and Remark 1.

(ii) Assume that Algorithm 1 has a solution $t_1, ..., t_q$. Then this solution satisfies (2.3)-(2.4).

To see this, let $\tau'$ be the $\tau$ from Step 3 in the loop where $t_i$ was picked out from $S$. We now use $\tau$ to denote the final product of Algorithm 1. (A.7) will remain valid, and integrating gives

$$\bar{E}(K_i - \bar{S}_{\tau'})^+ = P_{i}^{K_i,*}. $$
which is what we needed to show.

(iii) Assume that (2.3)-(2.4) has a solution \( t_1, ..., t_q \). Then this solution coincides with the output of Algorithm 1.

To see this, consider the order statistic \( t_{(1)} \leq ... \leq t_{(q)} \). Obviously, if \( t_{(1)} = t_i \) for some \( i \), then this index \( i \) must be picked out in the first passage through the loop of Algorithm 1, and \( t_i \) must have the same value as the \( t_i \) picked out by the Algorithm. (If there are several \( t_i \)'s with the same smallest value, the same applies). Similarly, by induction, one assumes that \( t_{(1)} \leq ... \leq t_{(j)} \) coincide with the results of Algorithm 1, and it is then easy to see that \( t_{(j+1)} \) also coincides with the results of the Algorithm, in view of Proposition 4. Note that the termination condition in Step 0 cannot be triggered, since with the \( \tau \) from the previous step,

\[
\tilde{E}(K_i - \tilde{S}_{\tau \land t_i})^+ < \tilde{E}(K_i - \tilde{S}_\tau)^+.
\]

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